The exercises come from Ronald de Wolf's lecture notes dW19, Chapters 13].

## Exercises

1.) (H) dW19, Exercise 13.1]: The following problem is a decision version of the factoring problem:

Given positive integers $N$ and $k$, decide if $N$ has a prime factor $p \in\{k, \ldots, N-1\}$.
Show that if you can solve this decision problem efficiently (i.e., in time polynomial in the input length $n=\lceil\log N\rceil$ ), then you can also find the prime factors of $N$ efficiently.
2.) dW19, Exercise 13.3]: This exercise shows how to use BQP-algorithms as subroutines in other BQP-algorithms.
(a) $(\mathbf{H})$ Suppose $L$ is a language in BQP. Let $f$ be the corresponding Boolean function, so $f(x)=1$ iff $x \in L$. Show that there is a $w \leq \operatorname{poly}(n)$ and a polynomial-size quantum circuit $U$ that implements the following map for all $x \in\{0,1\}^{n}$ :

$$
\left|x, 0^{w+1}\right\rangle \mapsto \sqrt{p}|x, f(x)\rangle|\phi(x)\rangle+\sqrt{1-p}|x, 1-f(x)\rangle|\psi(x)\rangle
$$

where $p \geq 1-\exp (-n)$, and $|\phi(x)\rangle$ and $|\psi(x)\rangle$ are states of the $w$-qubit workspace.
(b) Show that there is a polynomial-size quantum circuit $V$ that (when restricted to the subspace where the workspace qubits are $|0\rangle$ ) is $\exp (-n)$-close in operator norm to the following unitary:

$$
O_{f}:\left|x, b, 0^{w}\right\rangle \mapsto\left|x, b \oplus f(x), 0^{w}\right\rangle
$$

for all $x \in\{0,1\}^{n}$ and $b \in\{0,1\}$.
(c) (H) Suppose $L$ is a language in BQP, and you have a polynomial-size quantum circuit for another language $L^{\prime}$ that uses queries to the language $L$ (i.e., applications of the unitary $O_{f}$ ). Show that the language $L^{\prime}$ is also in BQP: there is a polynomial-size quantum circuit for $L^{\prime}$ that doesn't need queries to $L$.

## Hints

Exercise 1: Use binary search, running the algorithm with different choices of $k$ to "zoom in" on the largest prime factor.

Exercise 2, a: Use the Chernoff bound, e.g., as in the last item of dW19, Appendix B.2] to make the error probability exponentially small.

Exercise 2.c: Use the error analysis of homework Nr. 4 [WW19, Exercise 4.4].

## References

[dW19] Ronald de Wolf. Quantum computing: Lecture notes (version 5), 2019. arXiv: 1907.09415 v 5

