

# 2022 Quantum Computing Homework Nr. 7

András Gilyén

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Homework exercises – you can earn 10 points in total! The first two exercises worth 2 points each, while the last is worth 6 = 1+1+2+2. On the last page you can find some hints where indicated by **(H)**. Submission deadline: 6th December 14:00.

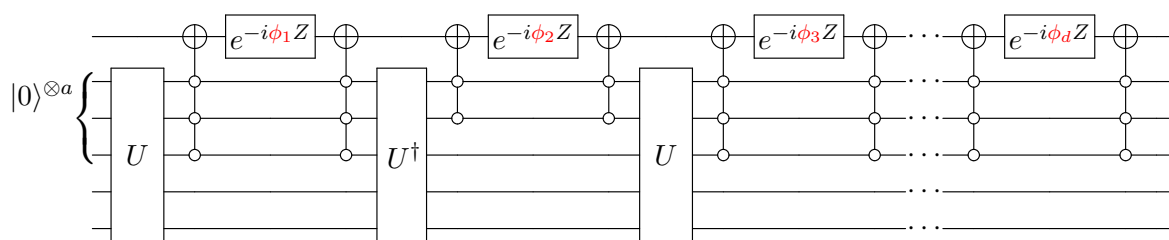
## Reminder

Let  $A = \sum_{i=1}^d \sigma_i |w_i\rangle\langle v_i|$  be a singular value decomposition of  $A$ , i.e.,  $v_i: i \in \{1, 2, \dots, d\}$  are orthonormal vectors as well as  $w_i: i \in \{1, 2, \dots, d\}$ , and the singular values  $0 \leq \sigma_i: i \in \{1, 2, \dots, d\}$  are ordered decreasingly.

**Theorem 1** (Quantum Singular Value Transformation (QSVT)[GSLW19]). *Let  $P: [-1, 1] \mapsto [-1, 1]$  be a degree- $d$  odd polynomial map. Suppose that  $U$  is a block-encoding of  $A = (\langle 0^a| \otimes I)U(|0^b\rangle \otimes I)$ . Then  $V := (H \otimes I)U_\Phi(H \otimes I)$  is a block-encoding of*

$$\sum_{i=1}^d P(\sigma_i) |w_i\rangle\langle v_i| = (\langle 0^{a+1}| \otimes I)V(|0^{b+1}\rangle \otimes I), \quad (1)$$

where  $\Phi \in \mathbb{R}^d$  is efficiently computable from  $P$  and  $U_\Phi$  is the following circuit.\*



## Exercises

- 1.) Show that the operator norm of  $A$  is  $\|A\| = \sigma_1$ .
- 2.) Let  $f: \mathbb{R} \rightarrow \mathbb{C}$ , and  $P \in \mathbb{C}[x]$  a polynomial such that  $|f(x) - P(x)| \leq \varepsilon$  for all  $x \in S \subseteq \mathbb{R}$ . Suppose that the singular values of  $A$  are elements of the set  $\sigma_i \in S$  for all  $i \in \{1, 2, \dots, d\}$ . Show that  $B := \sum_{i=1}^d f(\sigma_i) |w_i\rangle\langle v_i|$  and  $\tilde{B} := \sum_{i=1}^d P(\sigma_i) |w_i\rangle\langle v_i|$  are  $\varepsilon$ -close, i.e.,  $\|B - \tilde{B}\| \leq \varepsilon$ .
- 3.) This exercise shows how to solve a linear equation of the form  $Ax = b$ , when the equation might be under/over-determined. The least square solution is given by  $A^+b$ , where  $A^+$  is the Moore-Penrose pseudoinverse of  $A$  defined as  $A^+ := \sum_{i: \sigma_i \neq 0} \frac{1}{\sigma_i} |v_i\rangle\langle w_i|$ .

- (a) Prove that  $(AA^+)A = A$  and  $A^+(AA^+) = A^+$ . This shows that  $A^+$  indeed acts as we expect from a generalized inverse.

\*The empty dots denote control on the state  $|0\rangle$ . The generalized  $CNOT$ /Toffoli gates are controlled by  $|0^a\rangle$  and  $|0^b\rangle$  on the right- and left-hand sides of  $U$  respectively in the circuit – in this example circuit  $a = 3$ ,  $b = 2$ .

- (b) Suppose that  $U$  is a block-encoding of  $A$ , s.t.,  $A = (\langle 0| \otimes I)U(|0\rangle \otimes I)$ , and  $\|A^+\| \leq \kappa$ . Describe a bounded set  $S$  that contains the singular values of any such matrix  $A$ , but is disjoint from some interval of the form  $(0, a)$  for some  $a > 0$ . What is the largest  $a$  that we can choose?
- (c) Construct an approximate block-encoding  $V$  of  $A^+/(2\kappa)$  such that

$$\|A^+/(2\kappa) - (\langle 00| \otimes I)V(|00\rangle \otimes I)\| \leq \varepsilon,$$

and  $V$  uses  $\mathcal{O}(\kappa \log(1/\varepsilon))$ -times the block-encoding  $U$  and its inverse  $U^\dagger$ . (**H**)

- (d) Assume for simplicity that  $A^+/(2\kappa) = (\langle 00| \otimes I)V(|00\rangle \otimes I)$  holds exactly. Suppose that  $W$  is a quantum circuit that maps  $|0^n\rangle \rightarrow |b\rangle$ . Give a quantum algorithm that prepares a quantum state proportional to  $A^+|b\rangle$  with high probability with  $\mathcal{O}(\kappa)$  uses of the quantum circuits  $V$  and  $W$ . What does it tell us about the complexity of preparing a quantum state proportional to the least-square solution  $A^+|b\rangle$  given the block-encoding of  $A$ ?

## Hints

Exercise 3c: You can use the following polynomial approximation result [Gil19, Corollary 3.4.13.] without proof: for every  $\kappa > 1$  and  $\varepsilon < 1$  there is an odd polynomial  $P \in \mathbb{R}[x]$  of degree  $\mathcal{O}(\kappa \log(1/\varepsilon))$ , such that  $|P(x)| \leq 1$  for all  $x \in [-1, 1]$ , and  $|P(x) - 1/(2\kappa x)| \leq \varepsilon$  for all  $x \in [1/\kappa, 1]$ .

## References

- [Gil19] András Gilyén. *Quantum Singular Value Transformation & Its Algorithmic Applications*. PhD thesis, University of Amsterdam, 2019.
- [GSLW19] András Gilyén, Yuan Su, Guang Hao Low, and Nathan Wiebe. Quantum singular value transformation and beyond: Exponential improvements for quantum matrix arithmetics. In *Proceedings of the 51st ACM Symposium on the Theory of Computing (STOC)*, pages 193–204, 2019. arXiv: 1806.01838